### Martin-Löf Random Brownian Motion

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#### Brownian motion as a random function

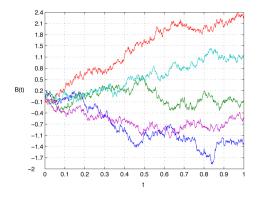


Fig. 1.1. Graphs of five sampled Brownian motions

Figure: Image from "Brownian Motion" by Mörters and Peres

Why Martin-Löf random Brownian Motion?

- Interesting random object
- More insight into "almost surely" results in classical theory of Brownian motion
- More insight into power of algorithmic randomness
- Recursion theoretic proofs of classical results

### Introduction

- Asarin and Prokovsky (1986): Complex Oscillations
- Willem Fouché (2000's): Strong foundations; proved Complex Oscillations are the Martin-Löf random sample paths of Brownian motion
- Laurent Bienvenu (2012): Use layerwise computability framework to do further work

## Technical Background

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### Classical Probability Theory

#### Definition (Classical Brownian Motion)

A collection of functions indexed by time ("real-valued stochastic process"),  $\{B(t) : t \in I\}$ , is called *standard Brownian motion* iff:

- Initial Value: B(0) = 0,
- Distribution: For all t ≥ 0 and h > 0, the increments B(t + h) B(t) are normally distributed with expectation 0 and variance h,
- Independent Increments: For all times  $0 \le t_1 \le t_2 \le \ldots \le t_n$ , the increments  $B(t_n) B(t_{n-1}), B(t_{n-1}) B(t_{n-2}), \ldots B(t_2) B(t_1)$  are independent random variables,
- Continuity: Almost surely, the function  $t \mapsto B(t)$  is continuous.

These requirements induce a measure on a function space, called *Wiener Measure*. We will talk about Martin-Löf randomness with respect to Wiener Measure on C[0, 1].

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### A little more technical detail

In particular, if  $A_1, ..., A_n$  are Borel sets in  $\mathbb{R}$ , then the probability of the finitary event  $[B(t_j) \in A_j \text{ for } 1 \le j \le n]$  is given by

$$\mathbb{P}[B(t_j) \in A_j \text{ for } 1 \le j \le n] =$$

$$\int_{A_1} \dots \int_{A_n} \prod \frac{1}{\sqrt{2\pi(t_j - t_{j-1})}} e^{\frac{-(v_j - y_{j-1})^2}{2(t_j - t_{j-1})}} dy_n \dots dy_1$$

### Franklin-Wiener Series Construction of Brownian Motion

We can construct Brownian motion from elements of Cantor space in the following way:

Each function will be represented by an infinite series of the form

$$B(t)=\xi_0\Delta_0(t)+\xi_1\Delta_1(t)+\sum_i\sum_{j<2^i}\xi_{i,j}\Delta_{i,j}(t)$$

where the  $\xi_{i,j}$  are real-number weights and the  $\Delta_{i,j}(t)$  are sawtooth functions.

### Franklin-Wiener Series Construction

- Using one binary real  $\alpha$ , we split into infinitely many binary reals  $\beta_0, \beta_1, \{\beta_{i,j}\}_{i \in \omega, j < 2^i}$
- We use a Gaussian distribution to map  $2^{\omega} \rightarrow [-\infty, \infty]$  and let  $\xi_{i,j} = g(\beta_{i,j})$

$$\beta = \frac{1}{\sqrt{2\pi}} \int_{-\infty}^{g(\beta)} e^{-t^2/2} dt$$

• See picture: Sawtooth functions  $\Delta_0, \Delta_1, \Delta_{i,j}$ 

### Franklin-Wiener Series Construction

$$B(t) = \xi_0 \Delta_0(t) + \xi_1 \Delta_1(t) + \sum_i \sum_{j < 2^i} \xi_{i,j} \Delta_{i,j}(t)$$

#### Theorem (Classical Probability Theory - due to Kahane?)

The Franklin-Wiener series representation converges almost surely to a continuous function in C[0,1], and the resulting class of functions is Brownian Motion.

#### Theorem (Fouché)

The Franklin-Wiener series representation converges to a continuous function for every  $\alpha \in 2^{\omega}$  that is Martin-Löf random.

### Layerwise computability

#### Definition (Layerwise Computable)

For  $U_n$  a universal Martin-Löf test on a space  $(X, \mu)$ , let  $K_n := X \setminus U_n$ . A function  $T : (X, \mu) \to Y$  is *layerwise computable* if it is computable on every  $K_n$ , uniformly on n.

#### Theorem (Hoyrup, Rojas)

If  $T : (X, \mu) \to Y$  is a layerwise computable map from a computable probability space to a computable metric space, then:

- The push-forward measure  $\nu \coloneqq \mu \circ T^{-1} \in \mathcal{M}(Y)$  is computable.
- T preserves Martin-Löf randomness; i.e. T(ML<sub>μ</sub>) ⊂ ML<sub>ν</sub>. Moreover, there is a constant c (computable from a description of T) such that T(K<sub>n</sub>) ⊂ K'<sub>n+c</sub> for all n, where K'<sub>n+c</sub> is the canonical layering of (Y, ν).

### Layerwise computability and MLR Brownian motion

#### Theorem (Fouché)

The Martin-Löf random paths of Brownian Motion are the image of Martin-Löf random reals in Cantor space under the Franklin-Wiener construction of Brownian Motion.

#### Theorem (Fouché)

Let B(t) be the Martin-Löf random Brownian motion constructed via the Franklin-Wiener series from a Martin-Löf random real  $\alpha$ . There is a piecewise linear function  $p_m$ , constructed from the first m bits of  $\alpha$ , such that

$$||B-p_m|| < \frac{\log m}{\sqrt{m}}$$

for all  $m > M_{\alpha}$ . Moreover,  $M_{\alpha}$  is layerwise computable from  $\alpha$ 

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## Useful Properties of Brownian Motion

### Classical Modulus of Continuity Results

#### Theorem (Levy)

For every constant  $c < \sqrt{2}$ , almost surely, for every  $\varepsilon > 0$  there exist  $0 < h < \varepsilon$  and  $t \in [0, 1 - h]$  with

$$|B(t+h) - B(t)| \ge c\sqrt{h\log(1/h)}$$

#### Theorem (Levy)

For every constant  $C > \sqrt{2}$ , almost surely, for every sufficiently small h > 0 and all 0 < t < 1 - h,

$$|B(t+h) - B(t)| \leq C\sqrt{h\log(1/h)}$$

### Modulus of Continuity of MLR Brownian motion

#### Theorem (Sharpening a result of Fouché)

Let B(t) be a Martin-Löf random Brownian motion. Then for all  $c < \sqrt{2}$ , for all  $h_0$ , there exists  $h < h_0$  such that for all t,

 $|B(t+h) - B(t)| \geq c\sqrt{h\log(1/h)}$ 

#### Theorem (Sharpening a result of Fouché)

Let B(t) be a Martin-Löf random Brownian motion. Then for all  $C > \sqrt{2}$ , there is an  $h_0$ , such that for all  $h < h_0$  and all t,

$$|B(t+h) - B(t)| \leq C\sqrt{h\log(1/h)}$$

and moreover,  $h_0$  is layerwise computable from B.

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## Zeros of Brownian Motion

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### Zeros of Brownian motion

#### Definition

For a path B(t) of one-dimensional Brownian motion, define  $Z_B = \{t \ge 0 : B(t) = 0\}$  to be the zero set of B

#### Theorem (Classical Probability Theory)

Almost surely,  $Z_B$  is a closed set with no isolated points.

#### Theorem

For B(t) a one-dimensional Martin-Löf random Brownian motion,  $Z_B$  is a closed set with no isolated points.

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### Characterize zeros of Martin-Löf Brownian motion

Necessary:

#### Theorem

If  $\alpha \in (0, 1]$  has effective Hausdorff dimension < 1/2,  $\alpha \notin Z_B$  for any Martin-Löf random Brownian motion B.

#### Sufficient:

#### Theorem

If  $\alpha \in (0, 1]$  has effective Hausdorff dimension > 1/2, then there exists a Martin-Löf random Brownian motion B such that  $\alpha \in Z_B$ .

Some points with effective dimension 1/2 are zeros, and some are not.

#### Theorem

For B a Martin-Löf random Brownian path, the first zero of B after any given computable real q is layerwise computable in B.

Proof:

#### Lemma (1)

It is layerwise computable in B(t) to see that there is a zero in a given interval [I, r] with computable endpoints.

#### Lemma (2)

It is layerwise computable in B(t) to see that there is not a zero in a given interval [I, r] with computable endpoints.

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### Proof of Lemma 1 - Demonstrating existence of a zero

- Main Idea: We can see, layerwise computably in *B*, positive and negative values in [*I*, *r*].
- (Fouché): The local minima (and maxima) of Martin-Löf random Brownian motion have non-recursive values.
- Thus 0 is never a local minimum or maximum
- So if a zero exists in [I, r] it is accompanied by intervals, arbitrarily close on either side, where B(t) > 0 and B(t) < 0.

### Proof of Lemma 1 - Demonstrating existence of a zero

- How do we determine, layerwise computably, if B(s) > 0 or B(s) < 0 for s computable?
- Recall (Fouché): From the first *m* bits of α ∈ 2<sup>ω</sup>, we can compute a piecewise linear function p<sub>m</sub>(t) such that

$$\|B(t)-p_m(t)\|\leq \frac{\log m}{\sqrt{m}}$$

for all  $m > M_{\alpha}$ 

- For a fixed s, we compute  $p_m(s)$  until  $||p_m(s)|| \ge \frac{\log m}{\sqrt{m}}$  and  $m > M_{\alpha}$ , which must eventually happen as  $B(s) \ne 0$  and  $\frac{\log m}{\sqrt{m}} \rightarrow 0$
- By running this computation on any dense set of computable reals, we must eventually observe positive and negative values for  $p_m(s)$  which are far enough from 0 that we know B(t) has crossed 0 in [l, r].

Proof of Lemma 2 - Seeing that there is not a zero

- Main Idea: We want to show, layerwise computably, that B(t) is bounded away from 0 on [I, r]
- Recall:  $\forall c > \sqrt{2} \exists h_0$  (layerwise computable in *B*) s.t.  $\forall h < h_0$

$$\|B(t+h) - B(t)\| < c\sqrt{h\log 1/h}$$

• So by sampling B(t) at intervals of size  $1/2^{n_i}$ ,  $n_i$  increasing and  $1/2^{n_0} < h_0$ , we eventually must find  $h^* = 1/2^{n_i}$  such that

$$B(t) > 2\sqrt{h^* \log 1/h^*}$$

for all  $t \in \{l + k/2^n\}_{k < 2^n}$  in [l, r].

### Completion of the proof

- Find the first zero of a Brownian path B(t) after a given computable real q
- Divide interval [q, 1] into intervals of size  $1/2^{n_0}$  for a suitable  $n_0$
- Run algorithms from lemmas to find the closest interval [*I*<sub>0</sub>, *r*<sub>0</sub>] to *q* that contains a zero.
- Divide  $[l_0, r_0]$  into intervals of size  $1/2^{n_1}$  and repeat
- Layerwise computably in B, we find sequences  $\{l_i\}$  and  $\{r_i\}$  which converge to the first zero after q from the left and from the right.
- Note: We could run a similar argument for crossing any computable real value after a computable time.

## An application to the Dirichlet Problem

### Brownian motion in other spaces

Brownian motion can be defined on  $C[0,\infty)$ , and can be constructed in a similar way by "gluing together" Brownian motion paths defined on [0,1]

#### Definition

 $B(t) = (B_1(t), B_2(t), ..., B_d(t)), t \ge 0$ , is *d*-dimensional (standard) Brownian motion if  $B_1, ..., B_d(t)$  are independent standard 1-dimensional Brownian motions.

#### Theorem (Fouché)

 $B(t) = (B_1(t), B_2(t), ..., B_d(t)), t \ge 0$ , is d-dimensional Martin-Löf random Brownian motion if  $B_1, ..., B_d(t)$  are mutually Martin-Löf random 1-dimensional Brownian motions.

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### **Dirichlet Problem**

#### Definition (Dirichlet Problem)

Given a function  $\phi$  that has values everywhere on the boundary of a region in  $\mathbb{R}^d$ , is there a unique continuous function u twice continuously differentiable in the interior and continuous on the boundary such that u is harmonic in the interior and  $u = \phi$  on the boundary?

### Dirichlet Problem

#### Theorem (Kakutani)

Suppose  $U \subset \mathbb{R}^d$  is a bounded domain such that every boundary point satisfies the Poincaré cone condition, and suppose  $\phi$  is a continuous function on the boundary  $\partial U$ . Let  $\tau(\partial U) = \inf\{t > 0 : B(t) \in \partial U\}$ , which is an almost surely finite stopping time. Then the function  $u : \overline{U} \to \mathbb{R}$  given by

$$u(x) = \mathbb{E}_x \left[ \phi(B(\tau(\partial U))) \right], \quad \text{ for } x \in \overline{U},$$

is the unique continuous function harmonic on U with  $u(x) = \phi(x)$  for all  $x \in \partial U$ .

### Computable Dirichlet Problem

#### Strong Conjecture

Let U be a bounded domain with computable boundary  $\partial U$ . Let  $\phi$  be a computable continuous function on  $\partial U$ . Then there is a unique, continuous, computable function  $u : \overline{U} \to \mathbb{R}$  harmonic on U such that  $u(x) = \phi(x)$  for all  $x \in \partial U$ .

### Further Directions

- Other types of randomness for Wiener measure
- Many "almost surely" results in classical theory
  - How much randomness do they require?
  - Do they characterize familiar notions of randomness?
- Can this more constructive approach be used to solve open problems in the classical theory?

# Thank You!

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