

# Density-one Points of $\Pi_1^0$ Classes

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Joe's and Noam's talks gave us an account of the class of density-one points restricted to the Martin-Löf random reals. Today we will extend this picture by saying a little about how they behave off the randoms.

We use the symbol  $\mu$  to refer exclusively to the standard Lebesgue measure on Cantor space.

Given  $\sigma \in 2^{<\omega}$  and a measurable set  $C \subseteq 2^\omega$ , the shorthand  $\mu_\sigma(C)$  denotes the relative measure of  $C$  in the cone above  $\sigma$ , i.e.,

$$\mu_\sigma(C) = \frac{\mu([\sigma] \cap C)}{\mu([\sigma])}.$$

### Definition

Let  $C$  be a measurable set and  $X$  a real. The **lower dyadic density** of  $C$  at  $X$ , written  $\rho_2(C|X)$ , is

$$\liminf_n \mu_{X \upharpoonright n}(C).$$

### Definition

A real  $X$  is a **dyadic positive density point** if for every  $\Pi_1^0$  class  $C$  containing  $X$ ,  $\rho_2(C|X) > 0$ . It is a **dyadic density-one point** if for every  $\Pi_1^0$  class  $C$  containing  $X$ ,  $\rho_2(C|X) = 1$ .

Our interest in density-one points originally developed in the setting of the real line. They arose in the study of effective versions of the Lebesgue Density Theorem, in the following form:

## Definition

Let  $C$  be a measurable subset of  $\mathbb{R}$  and  $x \in \mathbb{R}$ . The **lower (full) density** of  $C$  at  $x$ , written  $\rho(C|x)$ , is

$$\liminf_{\gamma, \delta \rightarrow 0^+} \frac{\mu((x - \gamma, x + \delta) \cap C)}{\gamma + \delta}.$$

## Definition

We say  $x \in [0, 1]$  is a **positive density point** if for every  $\Pi_1^0$  class  $C \subseteq [0, 1]$  containing  $x$ ,  $\rho(C|x) > 0$ . It is a **(full) density-one point** if for every  $\Pi_1^0$  class  $C \subseteq [0, 1]$  containing  $x$ ,  $\rho(C|x) = 1$ .

Recall two results from Joe's talk:

**Theorem (Bienvenu, Hölzl, Miller, Nies)**

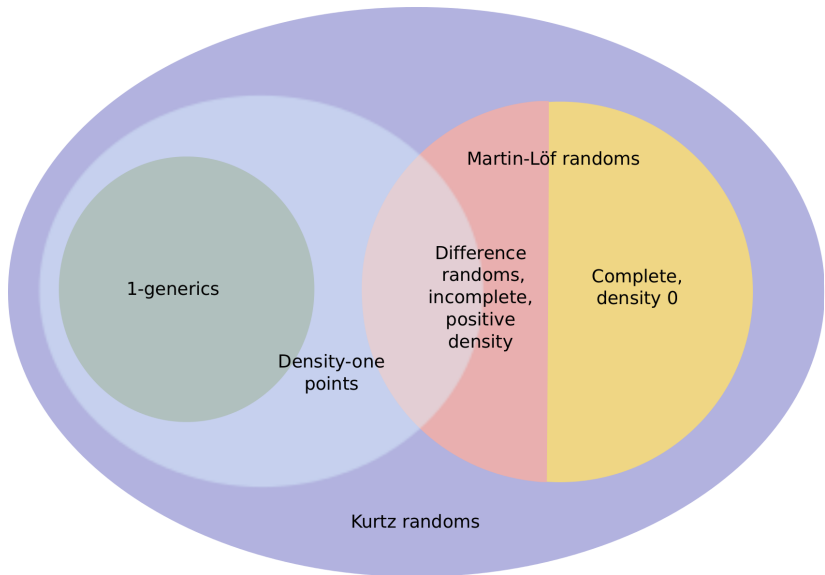
*If  $X$  is Martin-Löf random, then  $X$  is a positive density point if and only if it is incomplete.*

**Theorem (Day, Miller)**

*There is a Martin-Löf random real that is a positive density point (hence incomplete) but not a density-one point.*

- Dyadic positive density points (and hence full positive density points) are Kurtz random.
- 1-generics are full density-one points.
- Not being a full density-one point is a  $\Pi_2^0$  property. Therefore, all weak 2-random reals are full density-one points. Note that any hyperimmune-free Kurtz random is weak 2-random (Yu).
- The two halves of a dyadic density-one point are dyadic density-one. In fact, any computable sampling of a dyadic density-one point is a dyadic density-one point. Likewise for full density-one points.
- There is a Kurtz random real that is not Martin-Löf random and not a density-one point. Consider  $\Omega \oplus G$  where  $G$  is weakly 2-generic.

# The resulting picture



It's easy to exhibit a specific  $C$  and an  $X$  such that  $\rho_2(C | X) \neq \rho(C | X)$ . But is this discrepancy eliminated if we require that for every  $\Pi_1^0$  class  $C$  containing  $X$ ,  $\rho_2(C | X) = 1$ ? In other words, are dyadic density-one points the same as full density-one points? On the Martin-Löf randoms, yes:

### Theorem (K., Miller)

*Let  $X$  be Martin-Löf random. Then  $X$  is a dyadic density-one point if and only if it is a full density-one point.*

Some amount of randomness is necessary:

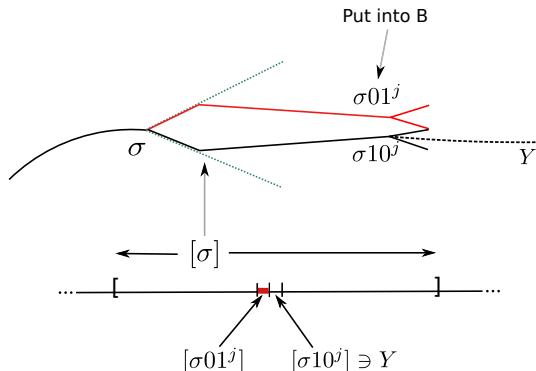
### Proposition (K.)

*There is a dyadic density-one point that is not full density-one.*

We build a dyadic density-one point  $Y$  by computable approximation, while building a  $\Sigma_1^0$  class  $B$  such that  $\rho(\bar{B} | Y) < 1$ .



The basic idea:



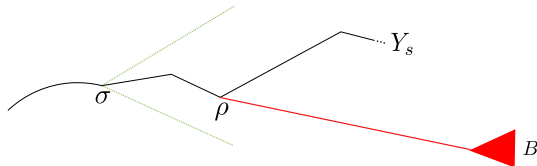
We shall be free to choose  $j$  as large as we want. Note that  $[\sigma]$  is the smallest dyadic cone containing  $Y$  that can see  $[\sigma 0 1^j]$ , the “hole” that we create in  $\bar{B}$ , and relative to  $\sigma$ , this hole appears small. However, on the real line, at a certain scale around  $Y$ , the hole is quite large.

## Dyadic density-one vs full density-one (contd.)

We want to place these holes infinitely often along  $Y$ , and this constitutes one type of requirement. Making  $Y$  a dyadic density-one point amounts to ensuring that for each  $\Sigma_1^0$  class  $[W_e]$ , either

- 1  $Y \in [W_e]$ , or
- 2 the relative measure of  $[W_e]$  along  $Y$  goes to 0.

The basic strategy for meeting a density requirement is to reroute  $Y$  to enter  $[W_e]$  if its measure becomes too big above some initial segment of  $Y_s$ . To make this play well with our hole-placing strategy, we keep the measure of  $B$  above initial segments of  $Y_s$  very small. Then if  $[W_e]$  becomes big enough, we can enter it while keeping  $B$  very small along  $Y_s$ .

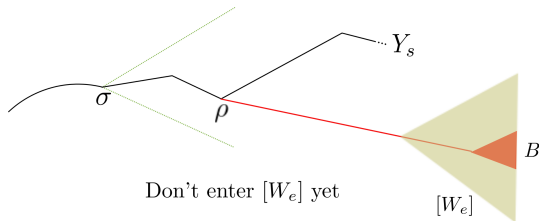


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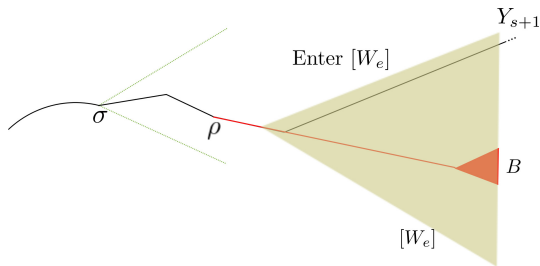


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The “Dyadic Density Drop Covering Lemma” makes this intuition precise:

## Lemma

Suppose  $B \subseteq 2^\omega$  is open. Then for any  $\varepsilon$  such that  $\mu(B) \leq \varepsilon \leq 1$ , let  $U_\varepsilon(B)$  denote the set

$$\{X \in 2^\omega : \mu_\rho(B) \geq \varepsilon \text{ for some } \rho \prec X\}.$$

Then  $U_\varepsilon(B)$  is open and  $\mu(U_\varepsilon(B)) \leq \mu(B)/\varepsilon$ .

(Joe proved this in his talk.)

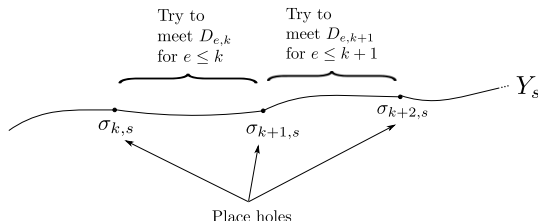
The lemma tells us exactly how small we have to keep  $B$  along  $Y_s$  to make it possible to act for multiple density requirements. Each time we reroute  $Y_s$  to enter a  $\Sigma_1^0$  class we get a little “closer” to  $B$ , but still remain far enough away so that we can act on behalf on another, higher priority density requirement if the need arises.

## Dyadic density-one vs full density-one (contd.)

Interleave hole-placing requirements with density requirements by progressively building a better and better approximation to a dyadic density-one point.

Formally, to meet the requirement  $D_{e,k}$  between  $\sigma$  and  $\sigma'$  where  $\sigma \preceq \sigma' \prec Y$  is to ensure that either  $\sigma' \in [W_e]$  or the measure of  $[W_e]$  between  $\sigma$  and  $\sigma'$  is bounded by  $2^{-k}$  (i.e., for every  $\tau$  between  $\sigma$  and  $\sigma'$ ,  $\mu_\tau([W_e]) \leq 2^{-k}$ ).

We organize the construction as follows:



$D_{e,k}$  has higher priority than  $D_{e',k}$ , for  $e' > e$ . Above  $\sigma_{k,s}$ , we only act for the sake of  $D_{e,k}$  if we haven't acted for the sake of a higher priority density requirement above  $\sigma_{k,s}$ . In sum, we have a finite-injury priority construction, where for each  $e$ , cofinitely many of the  $D_{e,k}$  requirements will be satisfied. There are some details to work out, but they're routine.  $\square$

1-generics are  $GL_1$ , therefore incomplete. By the theorem of Bienvenu et al., Martin-Löf random density-one points are also incomplete. But in general, density-one points can be complete. In fact, every real is computable from a full density-one point:

## Theorem (K.)

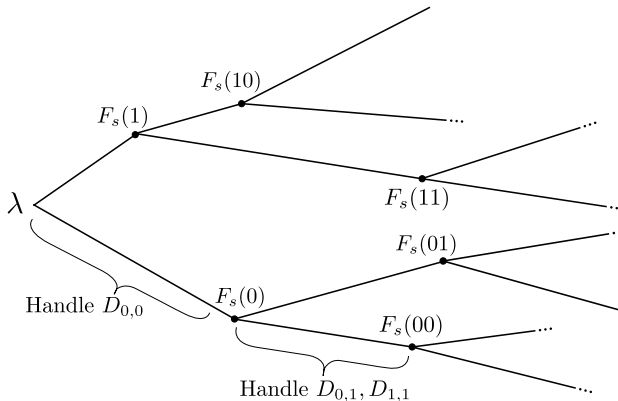
*For every  $X \in 2^\omega$ , there is a full density-one point  $Y$  such that  $X \leq_T Y \leq_T X \oplus 0'$ .*

Because dyadic density is so much easier to work with, I'll first sketch the proof of the result for dyadic density. Even though the statement of the theorem bears a superficial resemblance to the Kučera-Gács Theorem, the method is different. For one thing, there is no  $\Pi_1^0$  class consisting exclusively of density-one points. Also note that we don't get a  $wtt$  reduction as in the Kučera-Gács Theorem.

## Computational strength (contd.)

Basic idea: Combine the density strategy of the previous proof with coding, on a tree.

By computable approximation, we build a  $\Delta_2^0$  function tree  $F : 2^{<\omega} \rightarrow 2^{<\omega}$  and a functional  $\Gamma$  such that for every  $\sigma \in 2^{<\omega}$ ,  $\Gamma^{F(\sigma)} = \sigma$ .



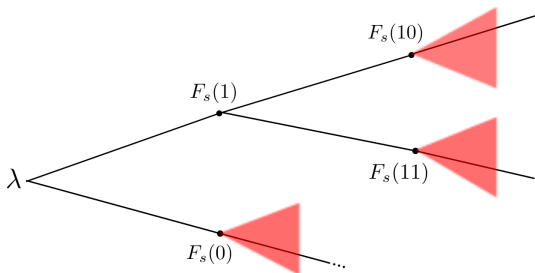


## Computational strength (contd.)

To set  $F_s(\sigma) = \tau$  at stage  $s$  is to code  $\sigma$  at the string  $\tau$ . We need to ensure that we can always do this in a consistent manner. There are two ways this could go wrong:

- $\tau$  codes incorrectly (i.e.,  $\Gamma^\tau \not\mid \sigma$ ), or
- $\tau$  codes too much (i.e.  $\Gamma^\tau$  properly extends  $\sigma$ ).

For example:



We cannot route  $F_s(1)$  through the current or previous values of  $F_s(10)$ ,  $F_s(11)$  and  $F_s(0)$ .

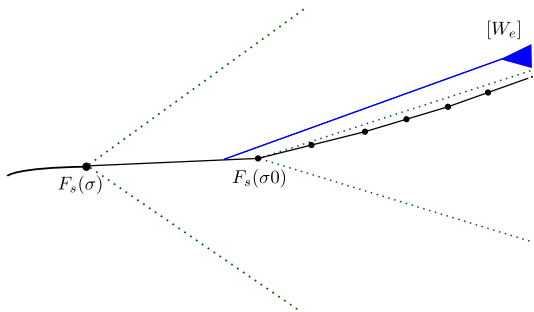
In general, for every nonempty string  $\sigma$ , there is a  $\Sigma_1^0$  class  $B_{\sigma,s}$  that  $F_s(\sigma)$  must avoid, and a threshold  $\beta_{\sigma,s}$  below which we must keep the measure of  $B_{\sigma,s}$  between  $F_s(\sigma^-)$  and  $F_s(\sigma)$ , where  $\sigma^-$  is the immediate predecessor of  $\sigma$ .

The strategies must cooperate to maintain this condition. For example, if  $\sigma = \alpha 0$ , then the strategies controlling  $F_s(\sigma 0)$ ,  $F_s(\sigma 1)$  and  $F_s(\alpha 1)$ , all of which contribute measure to  $B_{\sigma,s}$ , must maintain the fact that  $\mu(B_{\sigma,s})$  remains strictly below  $\beta_{\sigma,s}$  between  $F_s(\alpha)$  and  $F_s(\sigma)$ . All of this is completely within our control, since we can code on arbitrarily long strings.

For each  $X \in 2^\omega$ , the construction of  $\bigcup_{\sigma \prec X} F(\sigma)$  is again a finite-injury priority construction. The details are easy to work out. □

We briefly outline some of the difficulties in transferring the coding theorem for dyadic density-one points to full density-one points.

Strategies can no longer restrict their attention to dyadic cones:

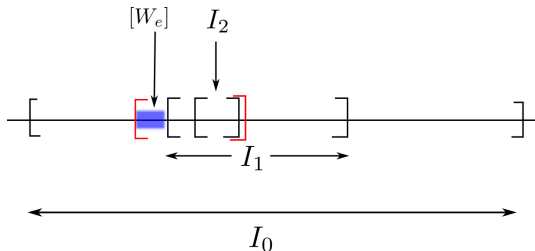


$[W_e]$  is very small relative to  $F_s(\sigma)$ , but it poses a threat to the path we're building.

We build a tree  $\{I_\sigma : \sigma \in 2^{<\omega}\}$  of intervals with dyadic rational endpoints.

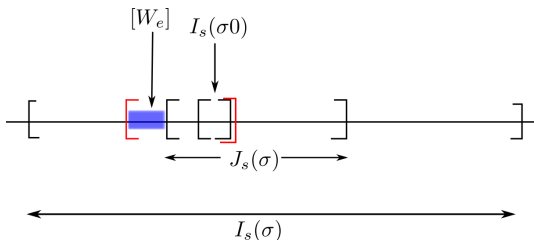
Suppose  $I \supseteq I'$  are intervals in  $[0, 1]$  and  $C$  is a measurable set. We say that  $\mu(C)$  is *below  $\varepsilon$  between  $I$  and  $I'$*  if for every interval  $L$  such that  $I \supseteq L \supseteq I'$ ,  $\mu_L(C) < \varepsilon$ .

In previous proofs, it was easy to chop up density requirements into smaller pieces such that the individual wins added up nicely. This is a little messier on the real line:

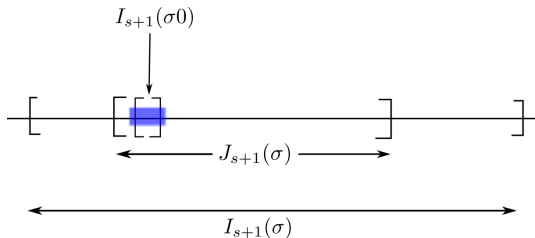


Here  $\mu([W_e]) < 1/8$  between  $I_0$  and  $I_1$  and also between  $I_1$  and  $I_2$ , but not between  $I_0$  and  $I_2$ .

So we watch for density drops on overlapping intervals. At every stage  $s$ , we maintain a proper subinterval  $J_s(\sigma)$  of  $I_s(\sigma)$  within which  $I_s(\sigma_0)$  and  $I_s(\sigma_1)$  reside. When the strategy in control of one of these intervals, say  $I_s(\sigma_0)$ , acts, it is allowed to move  $I_s(\sigma_0)$  outside  $J_s(\sigma)$ , in which case we expand  $J_s(\sigma)$  to an interval  $J_{s+1}(\sigma)$  that contains the new interval  $I_{s+1}(\sigma_0)$ .



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There is a version of the Density Drop Covering Lemma for the real line:

**Lemma (Bienvenu, Hölzl, Miller, Nies)**

*Suppose  $B \subseteq [0, 1]$  is open. Then for any  $\varepsilon$  such that  $\mu(B) \leq \varepsilon \leq 1$ , let  $U_\varepsilon(B)$  denote the set*

$$\{X \in [0, 1] : \exists \text{ an interval } I, X \in I, \text{ and } \mu_I(B) \geq \varepsilon\}.$$

*Then  $\mu(U_\varepsilon(B)) \leq 2\mu(B)/\varepsilon$ .*

We have to be slightly careful when applying this lemma for our construction. When we relativize this lemma to an interval  $L$ , we obtain a bound for the measure of  $U_\varepsilon(B \cap L)$  within  $L$ , but in general, we are also concerned about the part of  $B$  that lies outside  $L$ . Fortunately, under the assumptions of the construction, we can obtain a bound for the measure threatened by all of  $B$ .

We skip the details. On to the next topic...

There is a Kurtz random real of minimal degree. For example, every hyperimmune degree contains a weakly 1-generic (hence Kurtz random) set, and there is a minimal hyperimmune degree.

### Question

Is there a density-one point of minimal degree?

Of course, 1-generics and 1-randoms cannot be minimal, since for any real  $A \oplus B$  with either property,  $A$  and  $B$  are Turing incomparable. One might hope to prove the same for density-one points.



However, we think we can modify the  $\Delta_2^0$  construction of a density-one point with building a pair of Turing reductions between the two halves of the real:

### Conjecture (K.)

There is a density-one point  $A \oplus B$  with  $A \equiv_T B$ .

One approach is to try to show that some kind of computational strength that is compatible with minimality suffices to compute a density-one point. Even though highness is insufficient to compute a 1-generic, perhaps highness, or even hyperimmunity, implies the ability to compute a density-one point.

Note that a minimal density-one point would have to be hyperimmune.

- Every density 0 Martin-Löf random real is complete. Is there an incomplete Kurtz random real that is also density 0?
- Non-Martin-Löf random density-one points are hyperimmune. Do such points compute 1-generics?
- 2-random reals compute 1-generics (Kurtz; Kautz). On the other hand, there is a DNC function relative to  $0'$  of minimal degree (K.). Does every DNC function relative to  $0'$  compute a density-one point?
- Can a density-one point wtt-compute  $0'$ ?
- Are weakly 1-generics density-one points?

Thanks!